

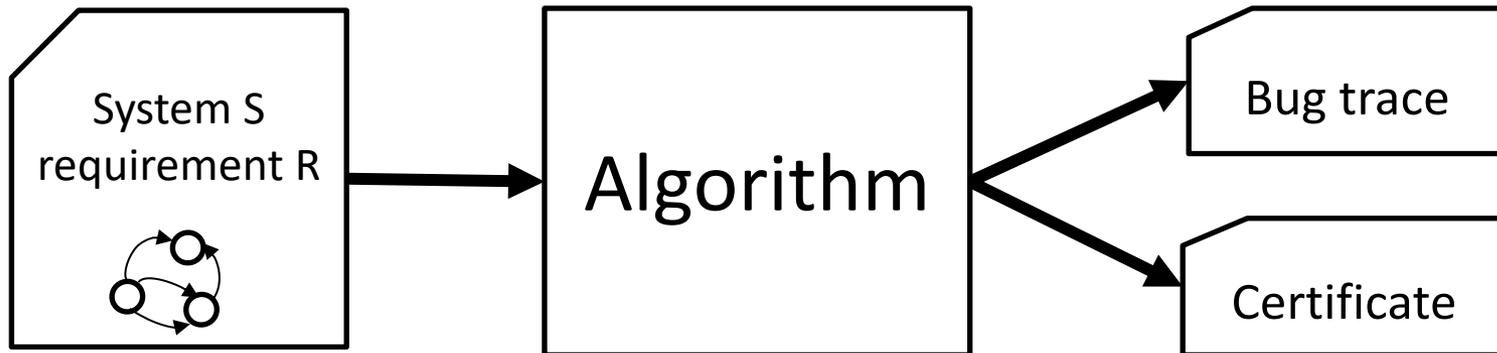
Data-driven verification

C2E2 and DryVR

Electrical & Computer Engineering
Coordinated Science Laboratory
University of Illinois at Urbana Champaign

ECE 584

Safety verification problem



Is there a behavior of system S violating safety requirement R within time bound T ?

Yes -> bug-trace -> design improvement

No -> safety proof -> certification

Safety verification problem

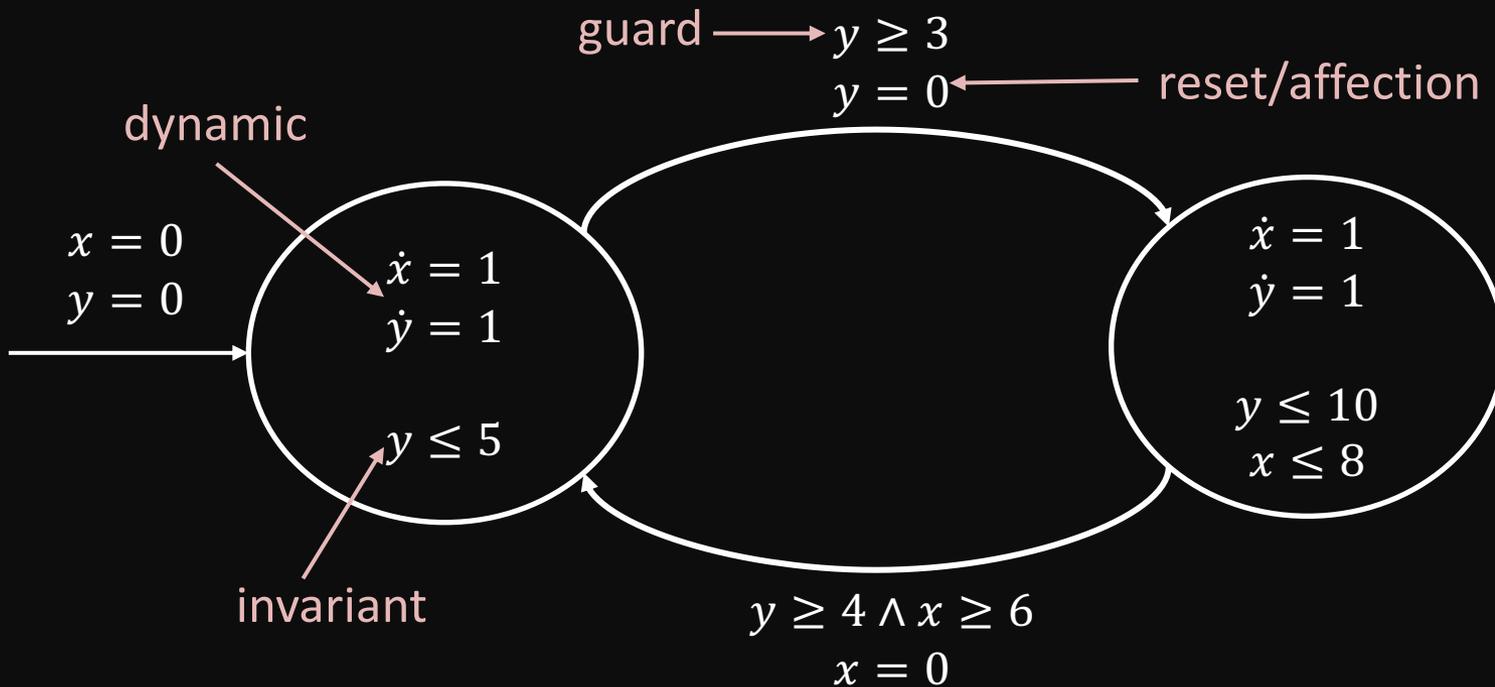


Is there a behavior of system S violating safety requirement R within time bound T ?

Yes \rightarrow bug-trace \rightarrow design improvement

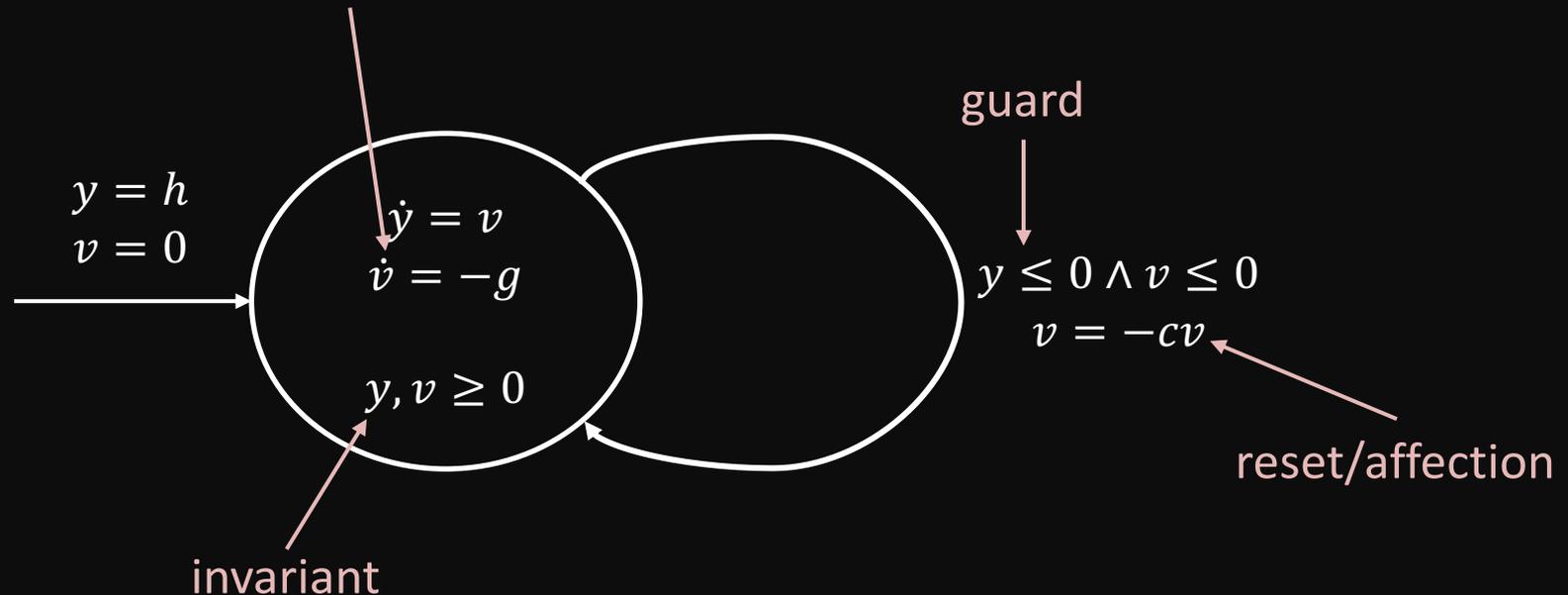
No \rightarrow safety proof \rightarrow certification

Recall: timed automata



Recall: bouncing ball

dynamic: general nonlinear function



Recall: bouncing ball



Avoid the Zeno behavior

Summary of C2E2

- Input: hyxml file
- Properties: initial set + unsafe set
- Simulate and/or verification
- Plotter

Outline

Introduction and C2E2 demo

Model-based sensitivity

- Simulation-driven verification algorithm
- Discrepancy function
- Matrix measure and sensitivity
- More examples

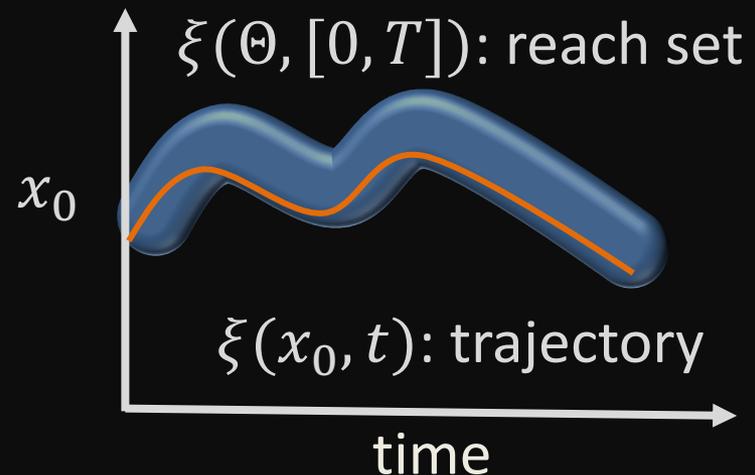
Next lecture on Thursday:

- New modeling questions with DryVR

System models and notations

nonlinear dynamical model

$$\dot{x}(t) = f(x(t))$$
$$\Theta, U \subseteq \mathbb{R}^n$$

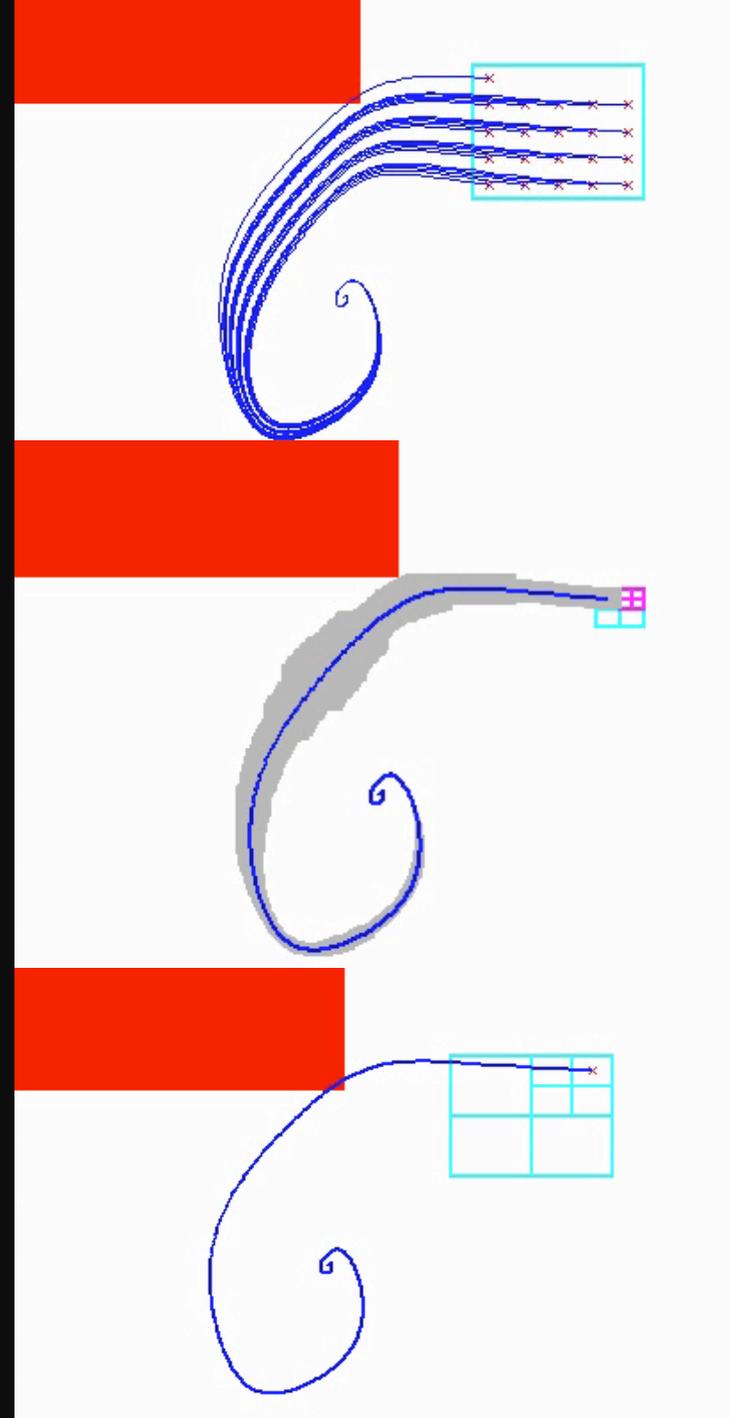


Safety verification problem $\xi(\Theta, [0, T]) \cap U = \emptyset?$

Simulations to safety proofs

- Given start Θ and target U
- Compute finite cover $\cup_i B(x_i, \delta) \supseteq \Theta$
- Simulate from the center x_0 of each cover to get $\xi(x_0, \{t_1, \dots, t_k\})$
- **Bloat** simulation so that
$$\xi(x_0, \cdot) \oplus \beta \supseteq \xi(B(x_0, \delta), [0, T])$$
- Check intersection/containment with U
- Refine cover if needed and repeat ...

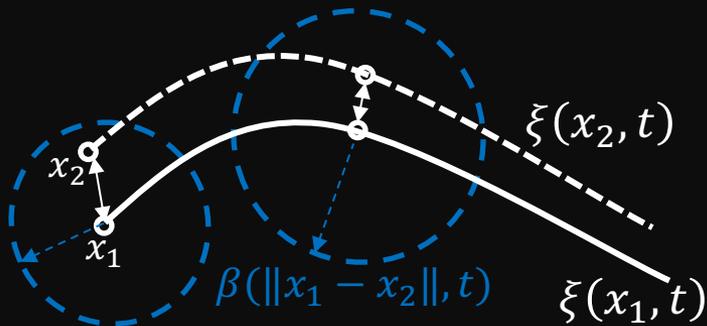
How to bloat or generalize simulations?



Brief history

2000	On Systematic Simulation of Open Continuous Systems	Kapinski et al.
2006	Verification using simulation	Girard and Pappas
2007	Robust Test Generation and Coverage for Hybrid Systems	Julius, Fainekos, et al.
2010	Breach, a toolbox for verification and parameter synthesis of hybrid systems.	Donzé
2013	Verification of annotated models from executions.	Duggirala, <i>Mitra</i> , Viswanathan

Main problem: How to quantify generalization?

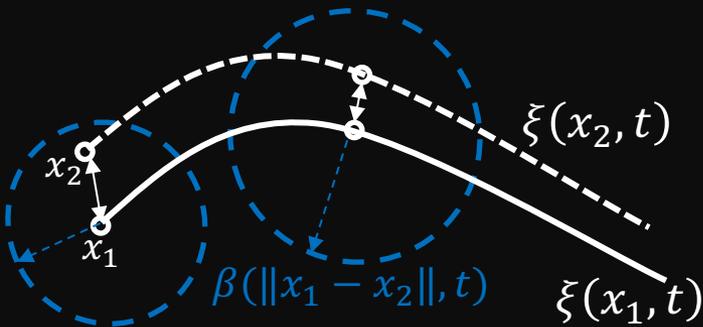


- Discrepancy formalizes generalization :
- Discrepancy is a continuous function β that bounds the distance between neighboring trajectories

$$\|\xi(x_1, t) - \xi(x_2, t)\| \leq \beta(\|x_1 - x_2\|, t),$$

- From a single simulation of $\xi(x_1, t)$ and discrepancy β we can over-approximate the reachtube

A simple example of discrepancy function



- If $f(x)$ has a Lipschitz constant L :

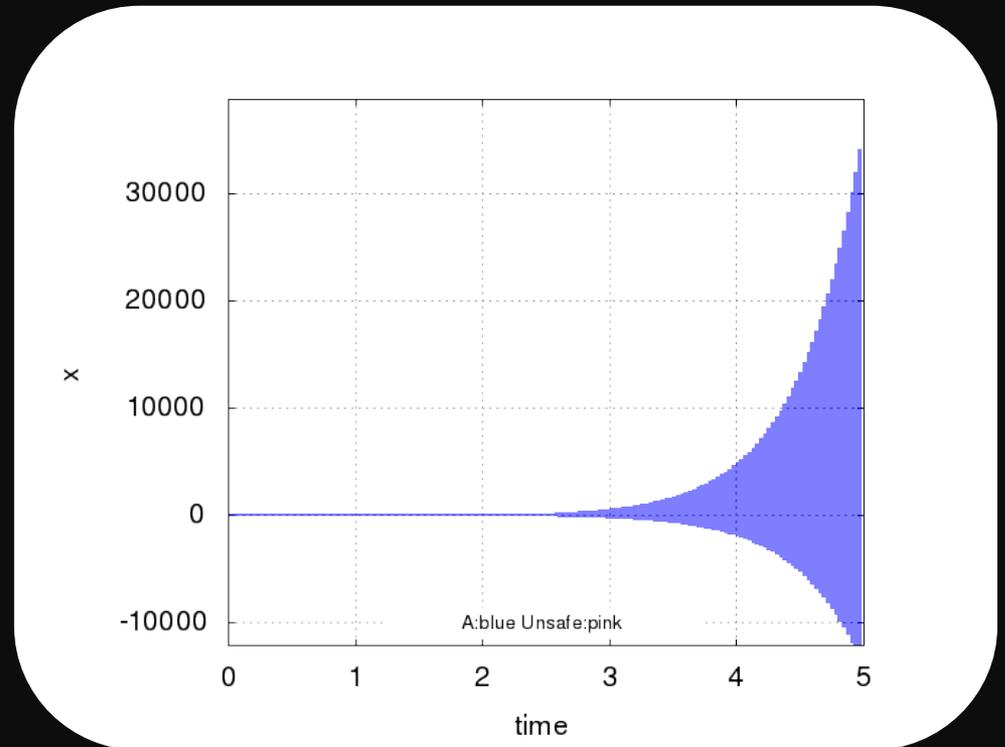
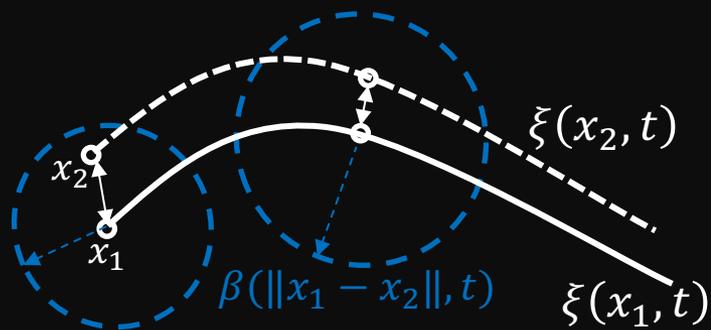
$$\forall x, y \in \mathbb{R}^n, \|f(x) - f(y)\| \leq L\|x - y\|$$

Example: $\dot{x} = -2x$, Lipschitz constant $L = 2$

- then a (bad) discrepancy function is

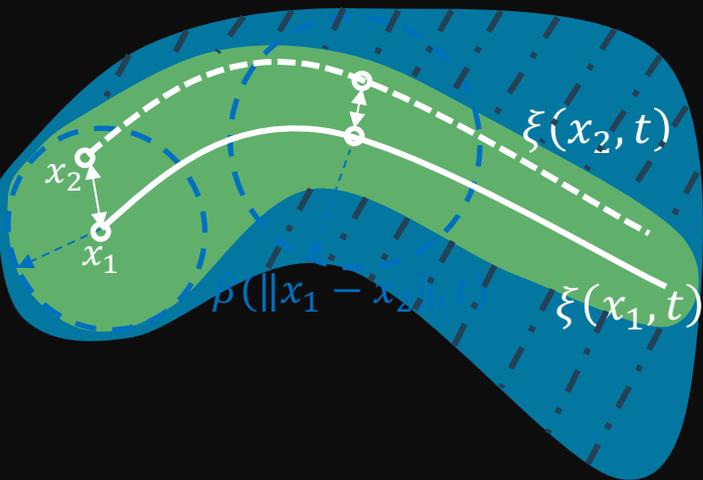
$$\|\xi(x_1, t) - \xi(x_2, t)\| \leq \|x_1 - x_2\|e^{Lt} = \beta(\|x_1 - x_2\|, t)$$

A simple example of discrepancy function



$$\dot{x} = -2x, \text{ Lipschitz constant } L = 2, \delta = 1$$

What is a good discrepancy ?



General: Applies to general nonlinear f

Accurate: Small error in β

Effective: Computing β is fast (in practice)

Matrix measure for $A \in \mathbb{R}^{n \times n}$

Matrix norm

$$\|A\| = \max_{x \neq 0} \frac{\|Ax\|}{\|x\|}$$

$$\|A\|_2 = \sqrt{\lambda_{\max}(A^T A)}$$

Matrix measure [Dahlquist 59]:

$$\mu(A) = \lim_{t \rightarrow 0^+} \frac{\|I + tA\| - \|I\|}{t}$$

$$\text{2-norm: } \mu(A) = \lambda_{\max}\left(\frac{A+A^T}{2}\right)$$

Computing μ

Vector norm	Induced matrix norm	Matrix measure
$ x _1 = \sum x_j $	$\ A\ _1 = \max_j \sum_i a_{ij} $	$\mu_1(A) = \max_j (a_{jj} + \sum_{i \neq j} a_{ij})$
$ x _2 = \sqrt{\sum x_j^2}$	$\ A\ _2 = \sqrt{\max_j \lambda_j(A^T A)}$	$\mu_2(A) = \max_j \frac{1}{2} (\lambda_j(A + A^T))$
$ x _\infty = \max_j x_j $	$\ A\ _\infty = \max_i \sum_j a_{ij} $	$\mu_\infty(A) = \max_i (a_{ii} + \sum_{j \neq i} a_{ij})$

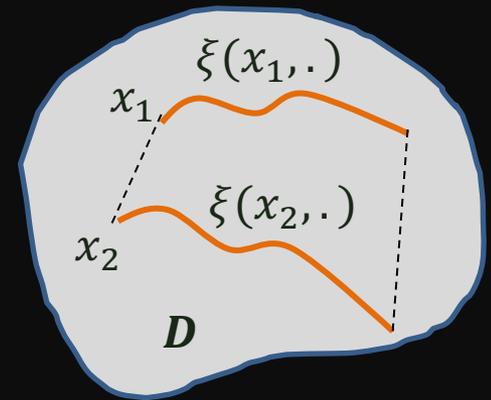
Table from: Reachability Analysis of Nonlinear Systems Using Matrix Measures [\[Maidens and Arcak, 2015\]](#)

Matrix measures can be used to compute discrepancy

Theorem [Sontag 10]: For any $\mathcal{D} \subseteq \mathbb{R}^n$, if the matrix measure of the Jacobian $\mu(J(t, x)) \leq c$ over \mathcal{D} , and all trajectories starting from the line remains in \mathcal{D} then the solutions satisfies:

$$|\xi(x_1, t) - \xi(x_2, t)| \leq |x_1 - x_2|e^{ct}$$

- That is, $|x_1 - x_2|e^{ct}$ is a discrepancy function
- Here J is the Jacobian of $f(x)$
- This c can be negative and is usually much smaller than the Lipschitz constant



Strategies for computing μ

- Define $y(t) = \xi(x_1, t) - \xi(x_2, t)$
- Let interval matrix \mathbf{A} be such that for all $x \in D, J_f(x) \in \mathbf{A}$,
- Then $\dot{y}(t) = A(t)y(t)$, for some $A(t) \in \mathbf{A}$
- This gives discrepancy $\beta \left(\|x_1 - x_2\|_M, t \right) = \|x_1 - x_2\|_M e^{\frac{\gamma^*}{2}t}$,
where $\gamma^* = \min \gamma$ s.t. $A^T M + M A \preceq \gamma M, \forall A \in \mathbf{A} \text{ --- } (*)$
- Solving (*)
 - Fix $M = I$, $\gamma^* = \lambda_{\max}(A + A^T) + \text{error}$

Simulation $\oplus \beta \rightarrow$ Reachtubes

simulation(x_0, h, ϵ, T) of gives sequence S_0, \dots, S_k :
 $\text{dia}(S_i) \leq \epsilon$ & at any time $t \in [ih, (i+1)h]$, solution
 $\xi(x_0, t) \in S_i$.

$\langle S_0, \dots, S_k, \epsilon_1 \rangle \leftarrow \text{valSim}(x_0, T, f)$

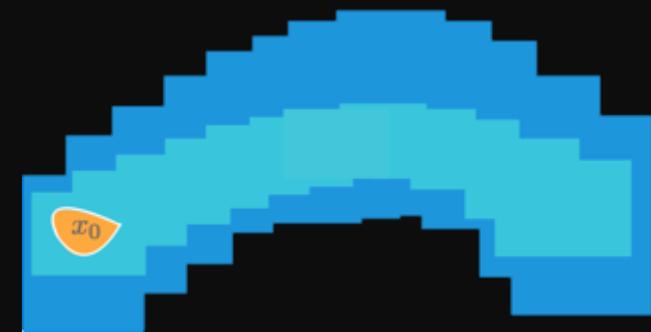
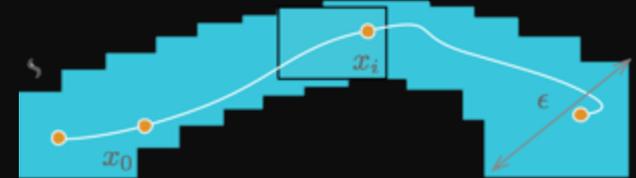
For each $i \in [k]$, $\epsilon_2 \leftarrow \sup_{t \in T_i, x, x' \in B_\delta(x_0)} \beta(x_1, x_2, t)$

$R_i \leftarrow B_{\epsilon_2}(S_i)$

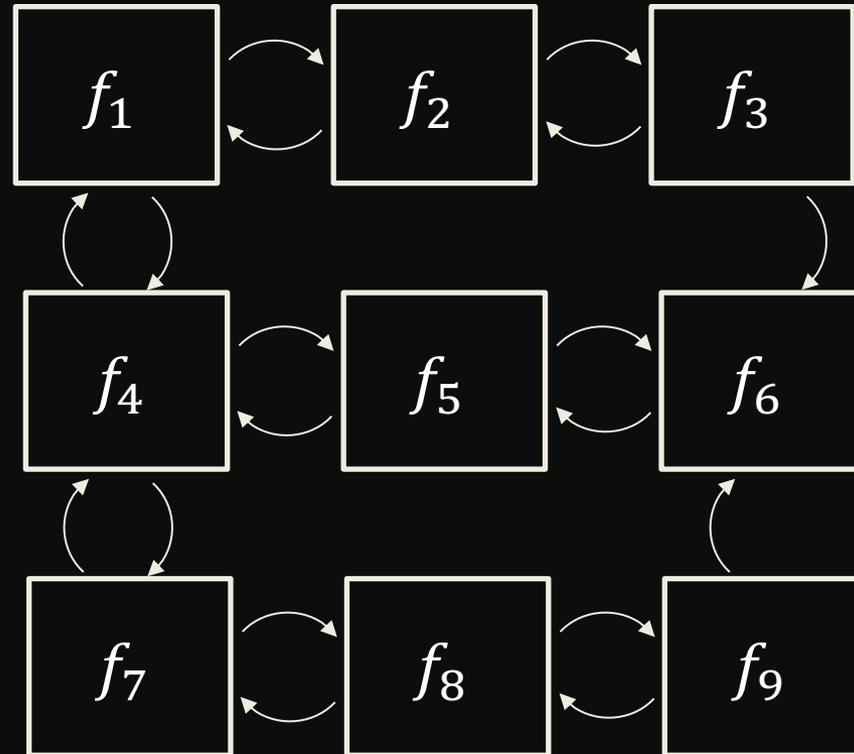
Example 1: $\dot{v} = \frac{1}{2}(v^2 + w^2); \dot{w} = -v$

- $J_f(v, w) = \begin{bmatrix} v & w \\ -1 & 0 \end{bmatrix}$
- $\gamma^* = 1.0178$ upper-bound on eigen values of the symmetric part of $J_f(v, w)$ over $D = [-2, -1] \times [2, 3]$
- $\|\xi(x_1, t) - \xi(x_2, t)\| \leq \|x_1 - x_2\| e^{1.0178t}$ while $x \in D$
- Uniform in all directions

Example 2: $\dot{x} = \begin{bmatrix} 0 & 3 \\ -1 & 0 \end{bmatrix} x$; Eigenvalues $\pm \sqrt{3} i$



Hybrid models



Hybrid Reachtubes

Track & propagate *may* and *must* fragments of reachtube

$$\mathit{tagRegion}(R, P) = \begin{cases} \mathit{must} & R \subseteq P \\ \mathit{may} & R \cap P \neq \emptyset \\ \mathit{not} & R \cap P = \emptyset \end{cases}$$

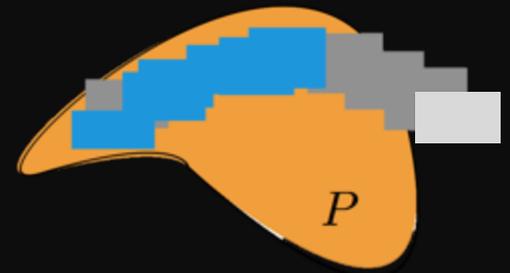
$\mathit{invariantPrefix}(\psi, S) =$

$\langle R_0, \mathit{tag}_0, \dots, R_m, \mathit{tag}_m \rangle$, such that either

$\mathit{tag}_i = \mathit{must}$ if all the R'_j s before it are *must*

$\mathit{tag}_i = \mathit{may}$ if all the R'_j s before it are at least *may*

and at least one of them is not *must*



Guarantees for bounded invariance verification using discreapancy

Theorem. (Soundness). If Algorithm returns safe or unsafe, then A is safe or unsafe.

Definition Given HA $A = \langle V, Loc, A, D, T \rangle$, an ϵ -perturbation of A is a new HA A' that is identical except, $\Theta' = B_\epsilon(\Theta)$, $\forall \ell \in Loc, Inv' = B_\epsilon(Inv)$ (b) $a \in A, Guard_a = B_\epsilon(Guard_a)$.

A is **robustly safe** iff $\exists \epsilon > 0$, such that A' is safe for U_ϵ upto time bound T , and transition bound N . Robustly unsafe iff $\exists \epsilon < 0$ such that A' is safe for U_ϵ .

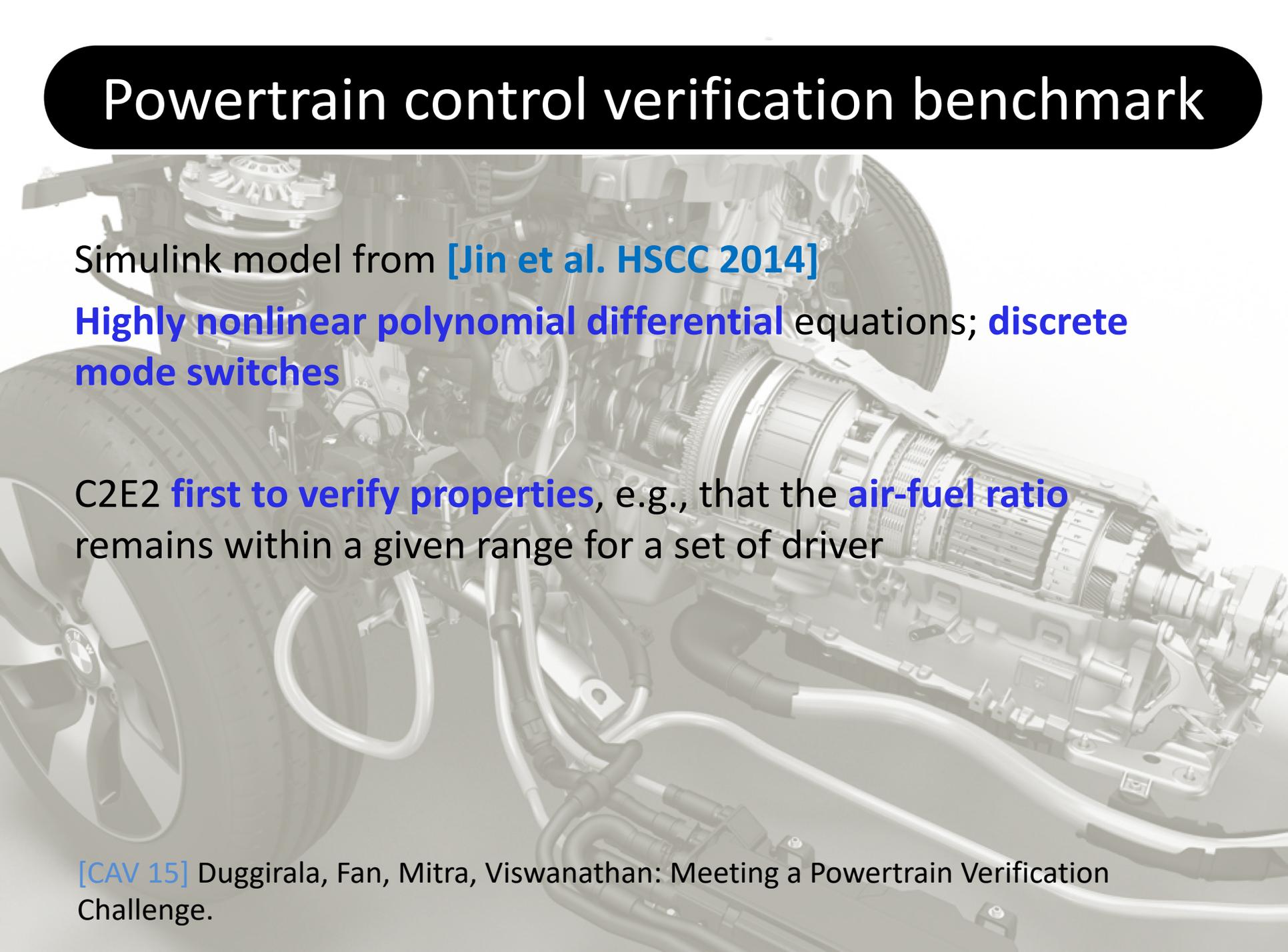
Theorem. (Relative Completeness) Algorithm always terminates whenever the A is either robustly safe or robustly unsafe.

Compare execute check engine



*static-dynamic analysis of
nonlinear hybrid models*

Powertrain control verification benchmark



Simulink model from [\[Jin et al. HSCC 2014\]](#)

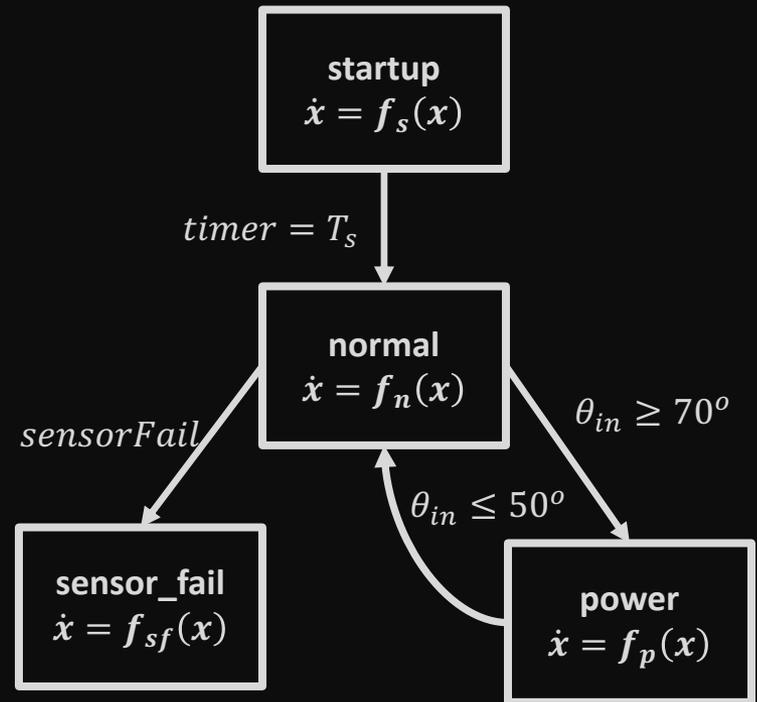
Highly nonlinear polynomial differential equations; discrete mode switches

C2E2 **first to verify properties**, e.g., that the **air-fuel ratio** remains within a given range for a set of driver

[\[CAV 15\]](#) Duggirala, Fan, Mitra, Viswanathan: Meeting a Powertrain Verification Challenge.

Polynomial hybrid automaton

Variable	Description
θ_{in}	Throttle angle
p	Intake manifold pressure
λ	Air/Fuel ratio
p_e	Intake manifold pressure estimate
i	Integrator state, control variable



$$\dot{\theta} = 10(\theta_{in} - \theta)$$

$$\dot{p} = c_1(2\theta(c_{20}p^2 + c_{21}p + c_{22}) - c_{12}(c_2 + c_3\omega p + c_4\omega p^2 + c_5\omega p^2))$$

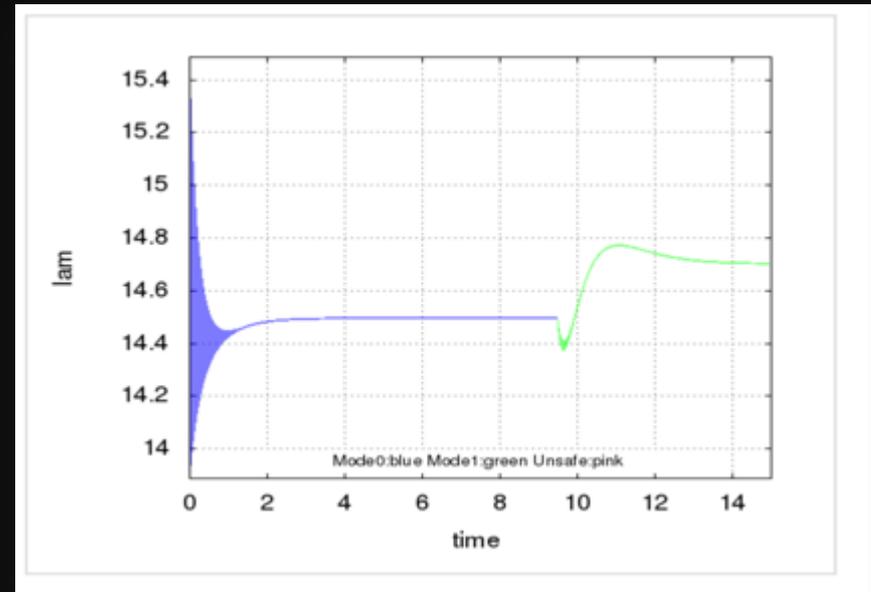
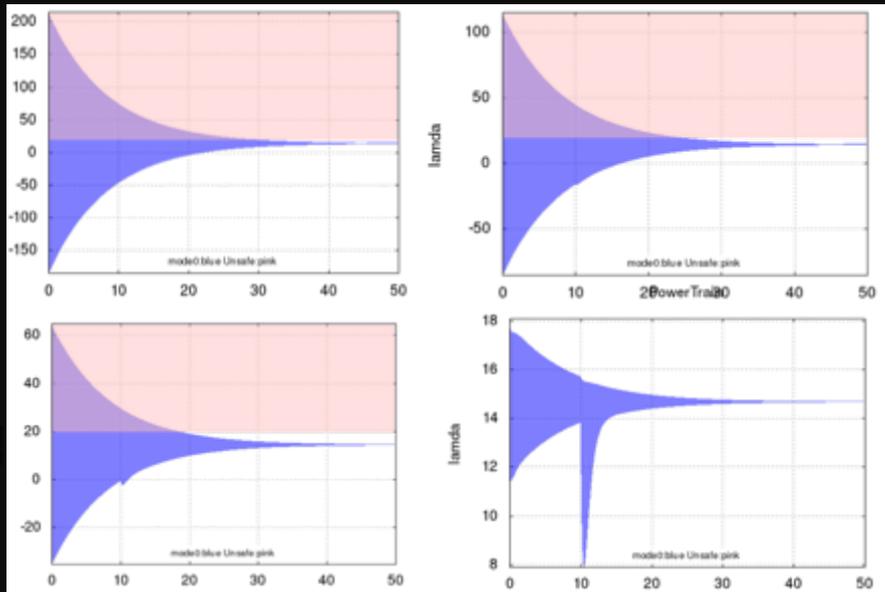
$$\dot{\lambda} = c_{26}(c_{15} + c_{16}c_{25}F_c + c_{17}c_{25}^2F_c^2 + c_{18}\dot{m}_c + c_{19}\dot{m}_c c_{25}F_c - \lambda)$$

$$\dot{p}_e = c_1(2c_{23}\theta(c_{20}p^2 + c_{21}p + c_{22}) - (c_2 + c_3\omega p + c_4\omega p^2 + c_5\omega p^2))$$

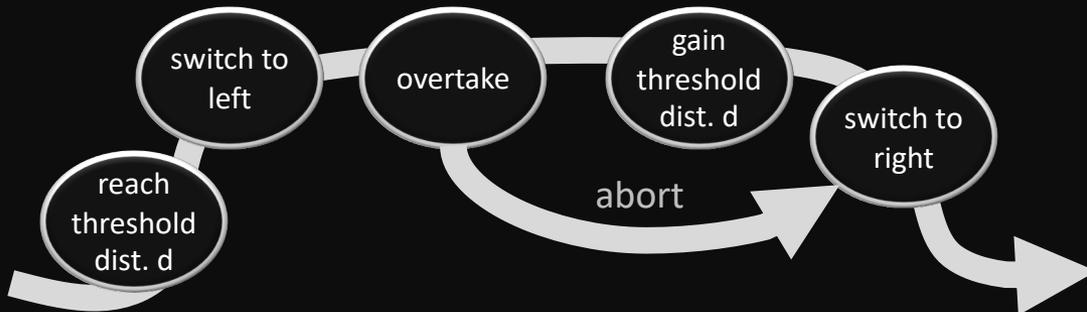
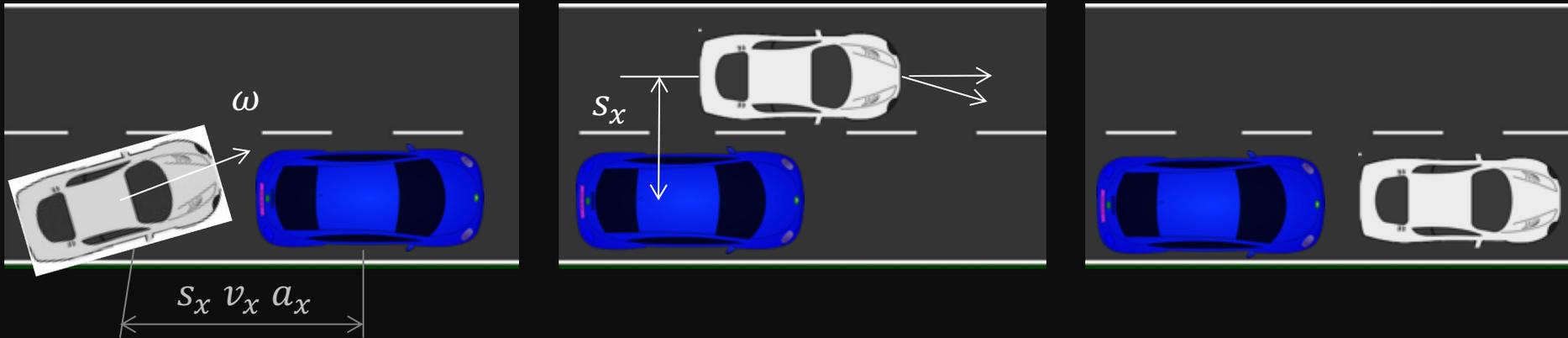
$$\dot{i} = c_{14}(c_{24}\lambda - c_{11})$$

Refinements in action: air-fuel ratio range

Requirement: Air-Fuel ratio λ contained in interval $[0.9\lambda_{ref}, 1.02\lambda_{ref}]$ for different initial conditions & throttle inputs



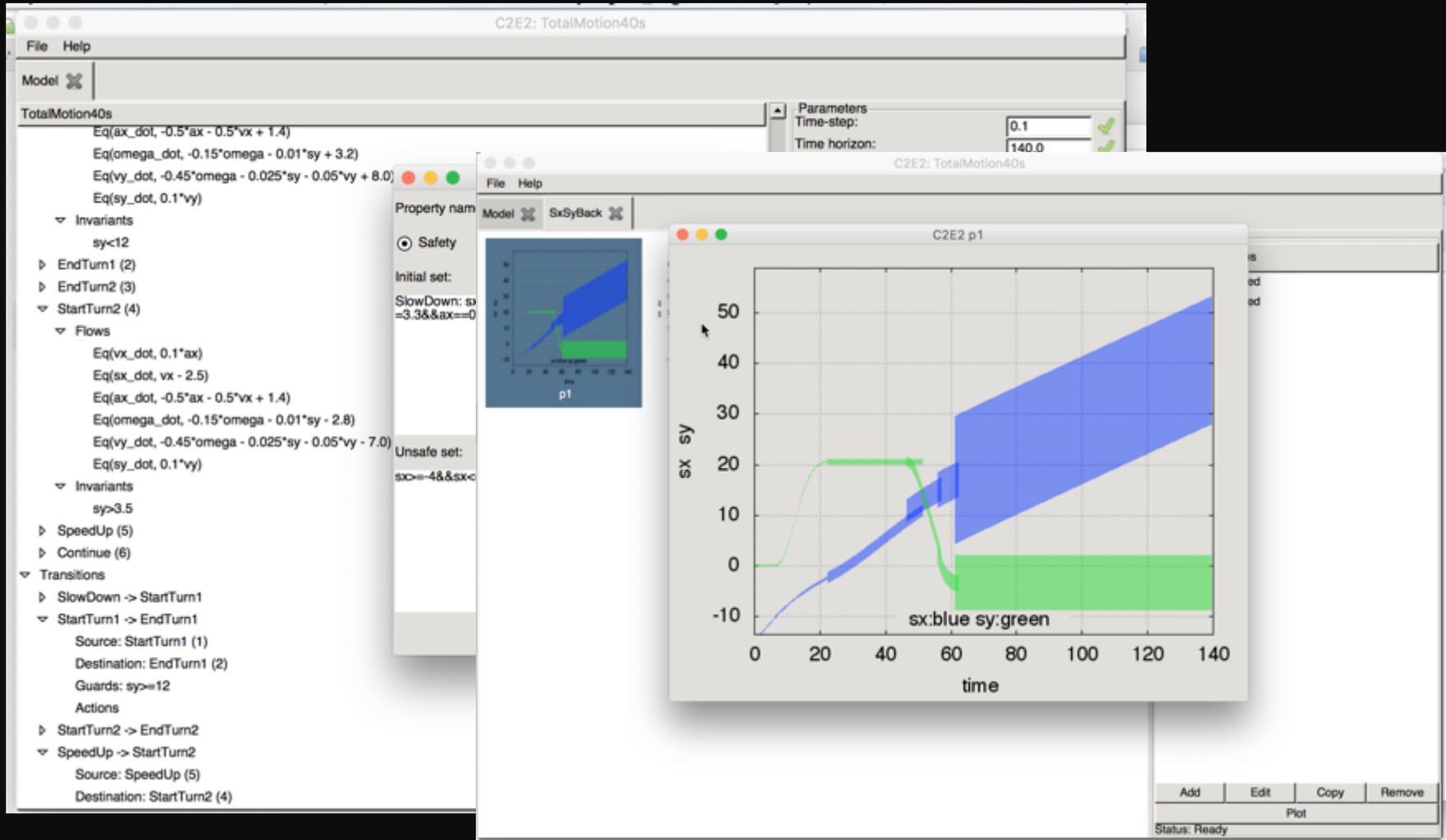
An auto-pass controller



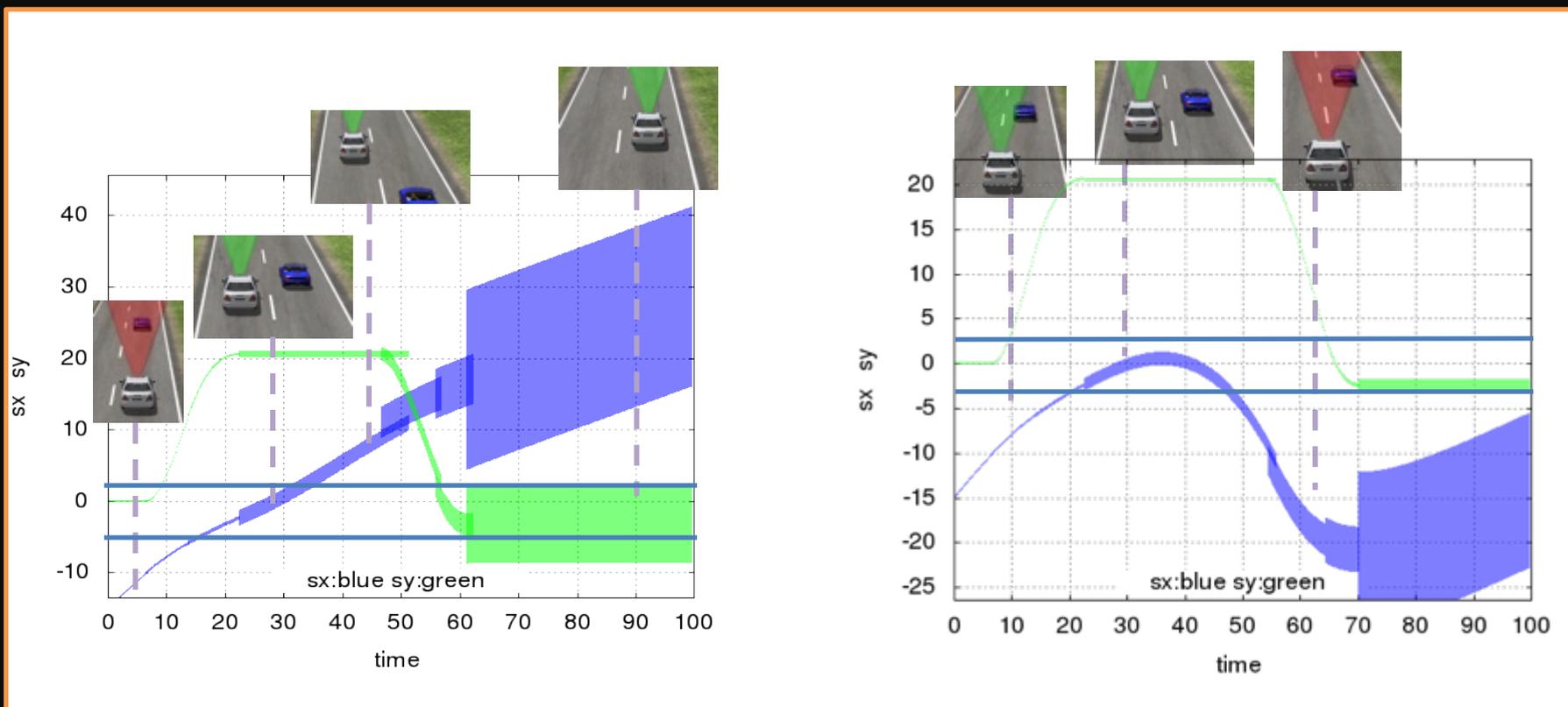
Given a controller and a safe separation requirement, we would like to check that the system is safe with respect to

- range of initial relative positions
- range of possible speeds
- range road friction conditions
- possible behaviors of "other" car
- range of design parameters

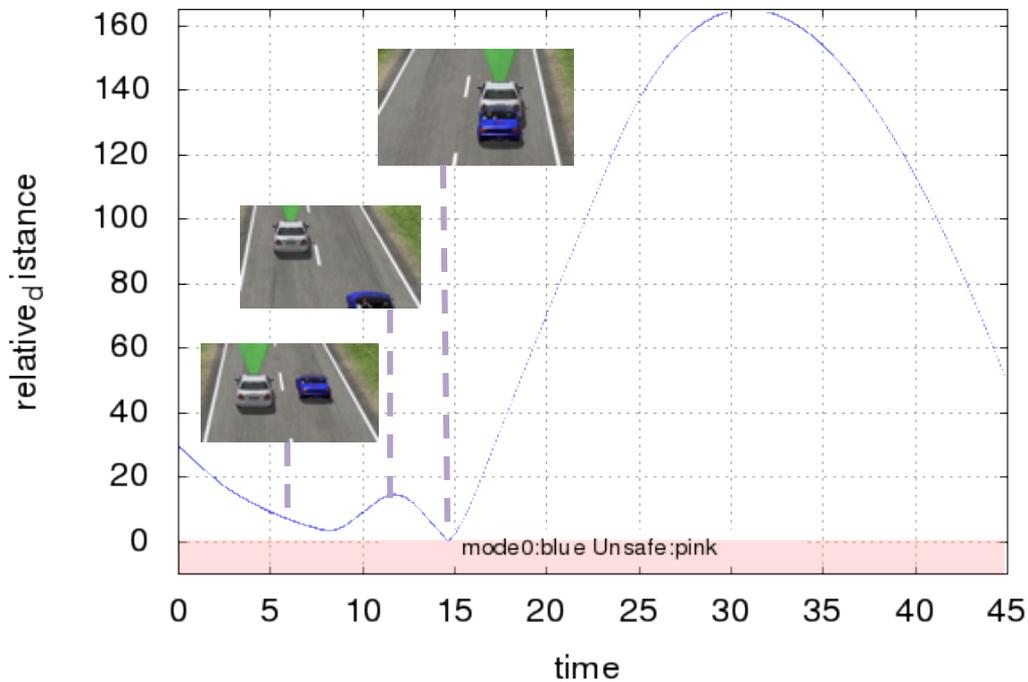
C2E2: Tool for nonlinear hybrid system verification



An auto-pass controller



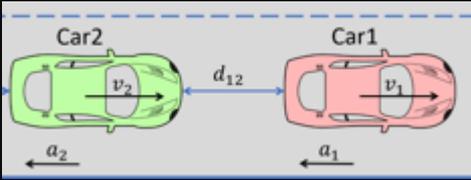
Debugging systems with high-fidelity models



Initial Set

Time Bound: 10s

Unsafe Set



Homework problem

Mode:Const_Const

Flow:

$$\begin{aligned}\dot{s}_1 &= v_1 \\ \dot{v}_1 &= 0 \\ \dot{s}_2 &= v_2 \\ \dot{v}_2 &= 0 \\ \dot{t} &= 1\end{aligned}$$

Inv: $t \leq 1$

Guard: $t \geq c_1$

Mode:Brake_Const

Flow:

$$\begin{aligned}\dot{s}_1 &= v_1 \\ \dot{v}_1 &= -2v_1 \\ \dot{s}_2 &= v_2 \\ \dot{v}_2 &= 0 \\ \dot{t} &= 1\end{aligned}$$

Inv: $s_1 - s_2 \geq 10$

Guard: $s_1 - s_2 \leq c_3$

Reset: $t = 0$

Mode:Brake_Brake

Flow:

$$\begin{aligned}\dot{s}_1 &= v_1 \\ \dot{v}_1 &= -2v_1 \\ \dot{s}_2 &= v_2 \\ \dot{v}_2 &= -3v_2 \\ \dot{t} &= 1\end{aligned}$$

Inv: $s_1 - s_2 \geq 0$

Reaction time

Guard: $t \geq c_2$

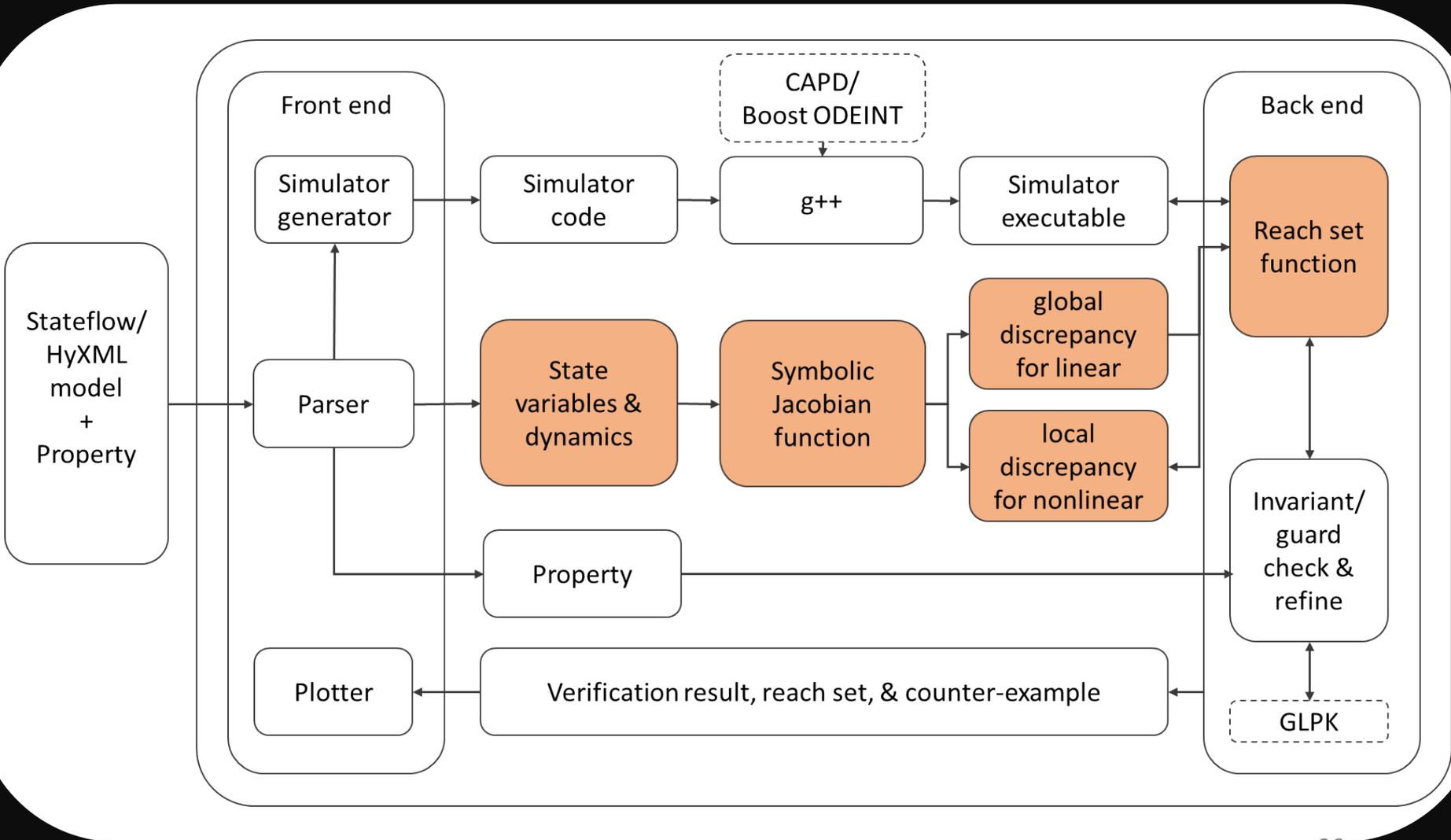
Mode:Brake_Const

Flow:

$$\begin{aligned}\dot{s}_1 &= v_1 \\ \dot{v}_1 &= -2v_1 \\ \dot{s}_2 &= v_2 \\ \dot{v}_2 &= 0 \\ \dot{t} &= 1\end{aligned}$$

Inv: $t \leq 0.4$

C2E2 Architecture



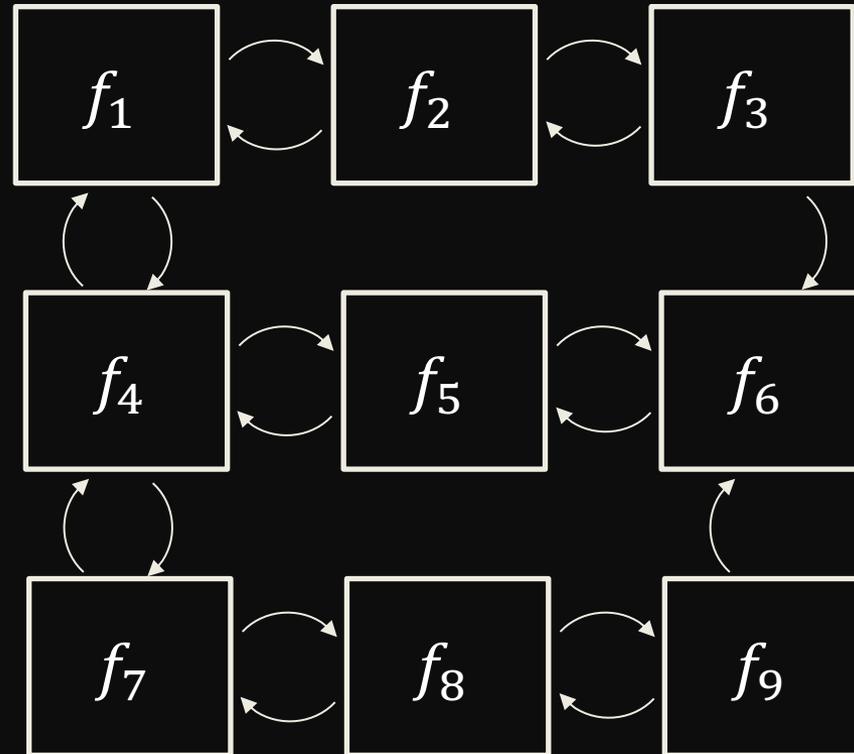
More features

- Log file to debug
- Plotted pictures are saved in the work-dir folder
- Command line version

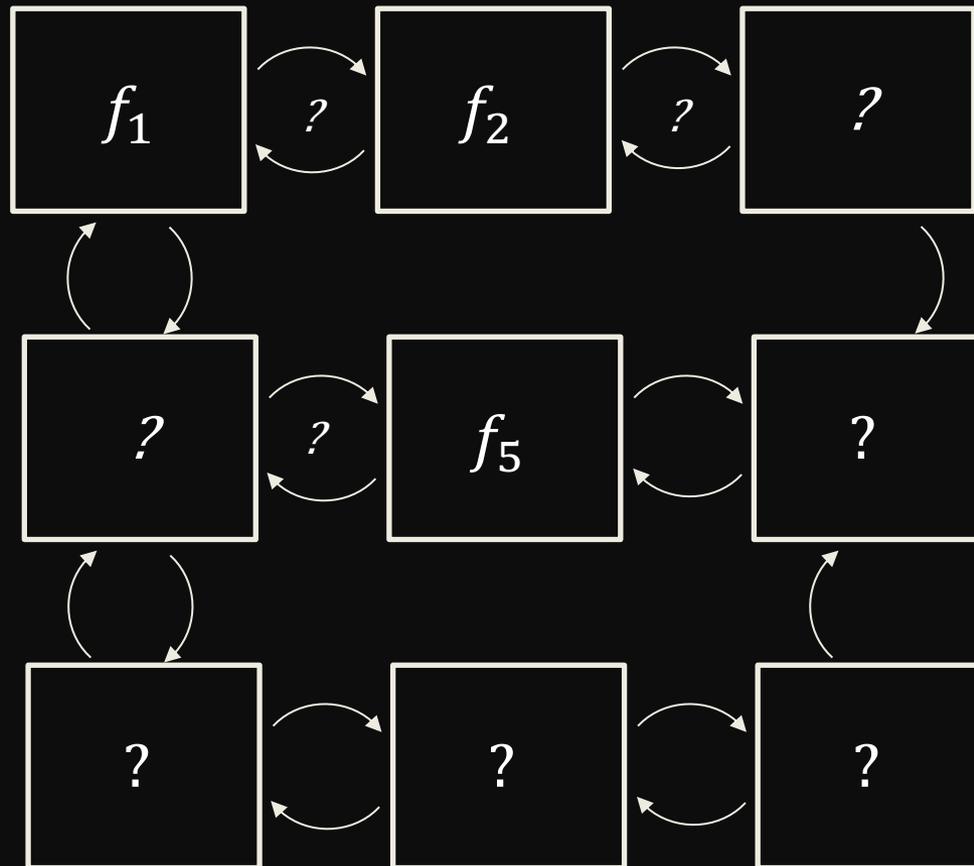
What we don't know

- Sample efficiency of the algorithms
 - Towards that [Girard Pappas 2006]
 - [Fan et al. EmSoft 2016] [Liberzon Mitra 2016]
- Unbounded initial set and time horizon
- How to verify open models?
 - $\dot{x}(t) = f(x(t), u(t)), x_0 \in \Theta, u \in \mathcal{U}$
 - Ongoing work with $\mathcal{U} = \{u_1, \dots, u_k\}$
- More general models with uncertainty

Hybrid models



Models closer to reality



“All models are wrong, some are useful”

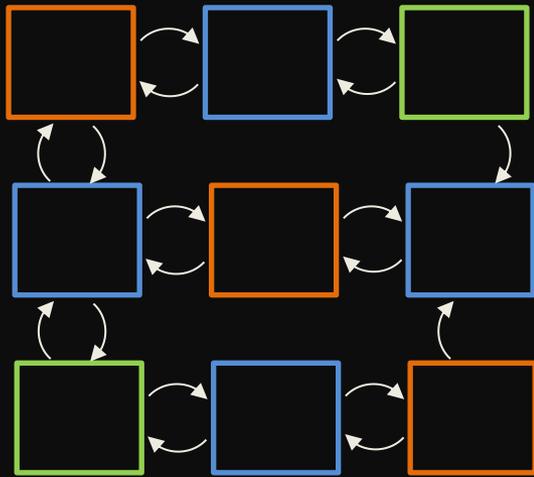
Dry  R

Gain serenity to accept models as they are

<https://github.com/qibolun/DryVR>

A new view of knowledge in hybrid models

Complete information
of switching structure



Transitions are time-
triggered, possibly
nondeterministic: one-
clock timed automaton

Executable access to
mode dynamics



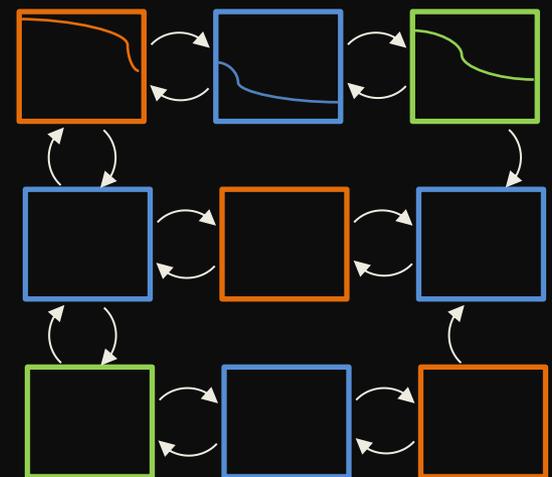
+



=

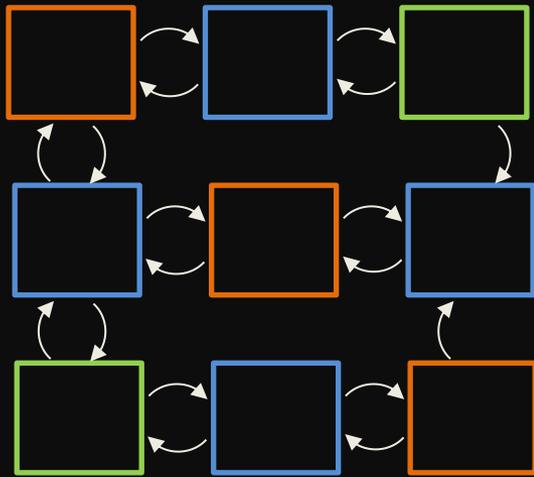


DryVR's Executable
hybrid model



A new view of knowledge in hybrid models

Formal reasoning
simulation, composition



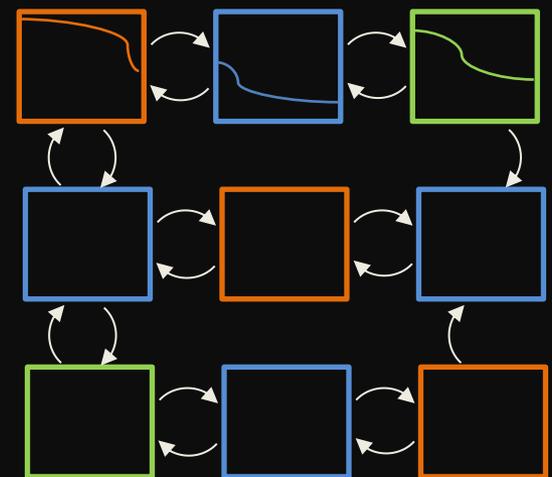
Statistical reasoning
sensitivity analysis



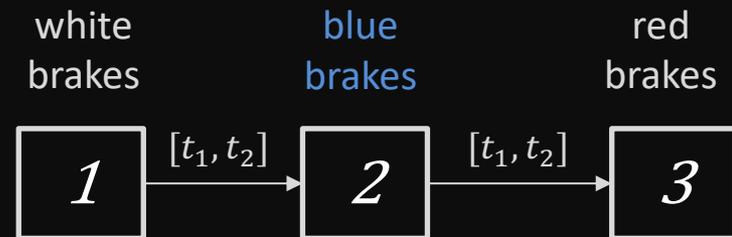
+

=

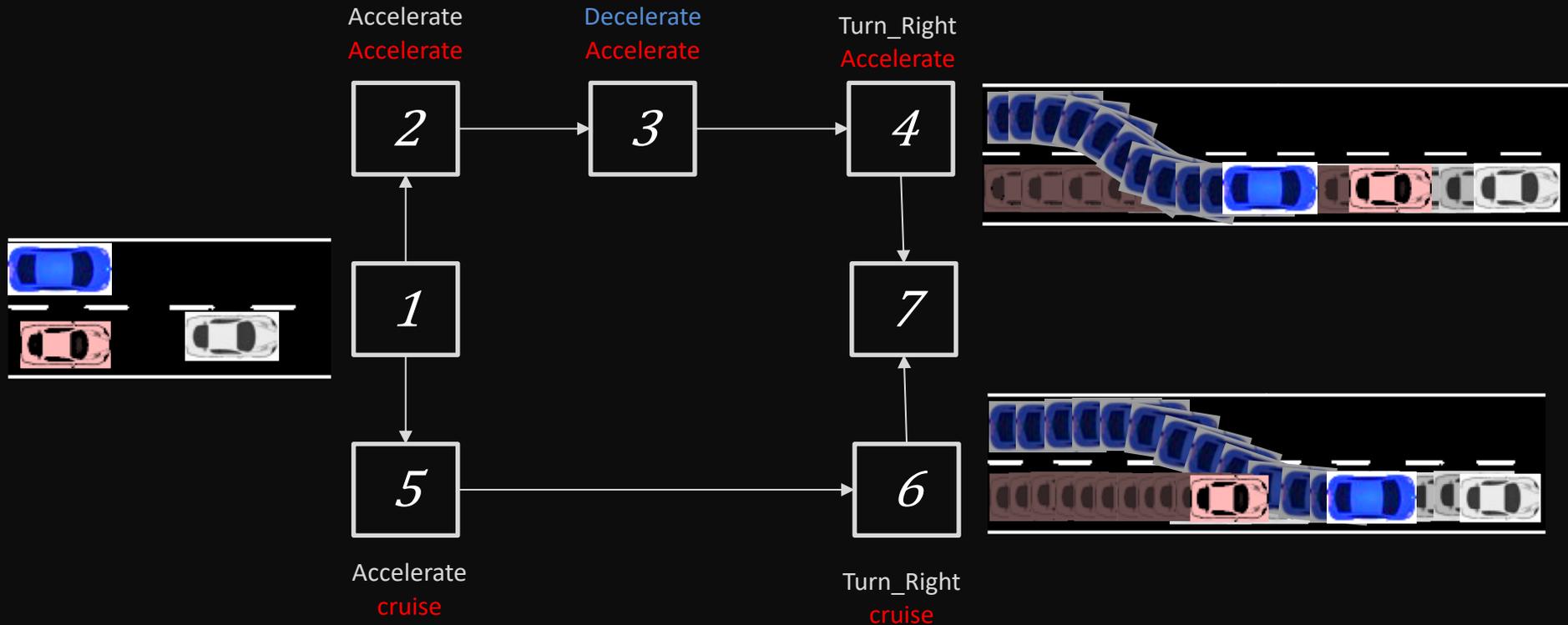
DryVR's formal
probabilistic guarantees



DryVR model for Automatic Emergency Breaking



DryVR model for auto-pass



Learning discrepancy from black-box

Assume a form of the discrepancy

Global exponential discrepancy

$$\beta(x_1, x_2, t) = |x_1 - x_2|K e^{\gamma t}$$

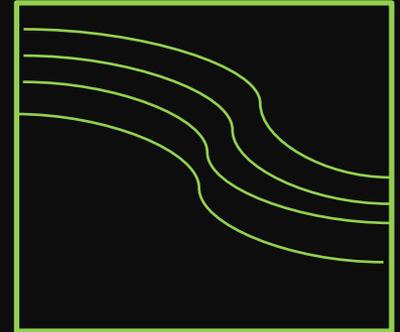
Others piece-wise exponential, polynomial

For any pair of trajectories τ_1 and τ_2 in mode \square

$$\forall t \in [0, T], |\tau_1(t) - \tau_2(t)| \leq |\tau_1(0) - \tau_2(0)|K e^{\gamma t}$$

$$\forall t, \ln \frac{|\tau_1(t) - \tau_2(t)|}{|\tau_1(0) - \tau_2(0)|} \leq \gamma t + \ln K$$

Familiar problem of learning linear separators



Learning linear separators

For a subset $\Gamma \subseteq \mathbb{R} \times \mathbb{R}$, a linear separator is a pair $(a, b) \in \mathbb{R}^2$ such that $\forall (x, y) \in \Gamma, x \leq ay + b$

Algorithm:

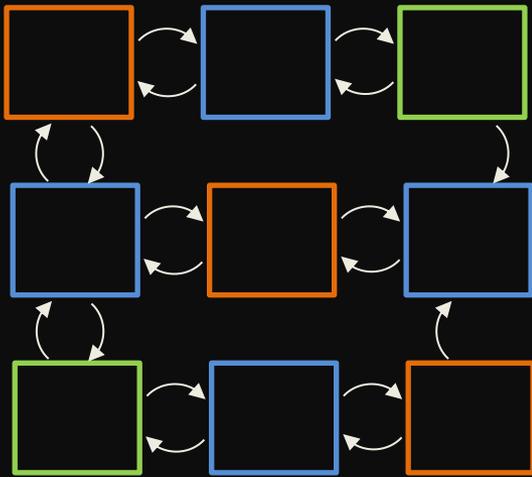
1. Draw k pairs $(x_1, y_1), \dots, (x_k, y_k)$ from Γ according to \mathcal{D} .
2. Find $(a, b) \in \mathbb{R}^2$ s.t. $x_i \leq ay_i + b$ for all $i \in \{1, \dots, k\}$.

Proposition [Valiant 84]: Let $\epsilon, \delta \in \mathbb{R}^+$. If $k \geq \frac{1}{\epsilon} \ln \frac{1}{\delta}$ then with probability $1 - \delta$, the above algorithm finds (a, b) such that $err_{\mathcal{D}}(a, b) = \mathcal{D}(\{(x, y) \in \Gamma \mid x > ay + b\}) < \epsilon$.

Experience: 96% accuracy for 10 trajectories, >99.9% for 20

DryVR

Complete information
of switching structure



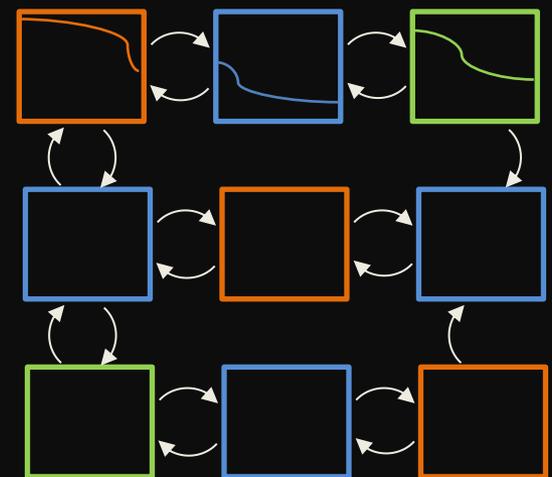
Model file specifies
vertices, edges, labels

Executable access to
mode dynamics

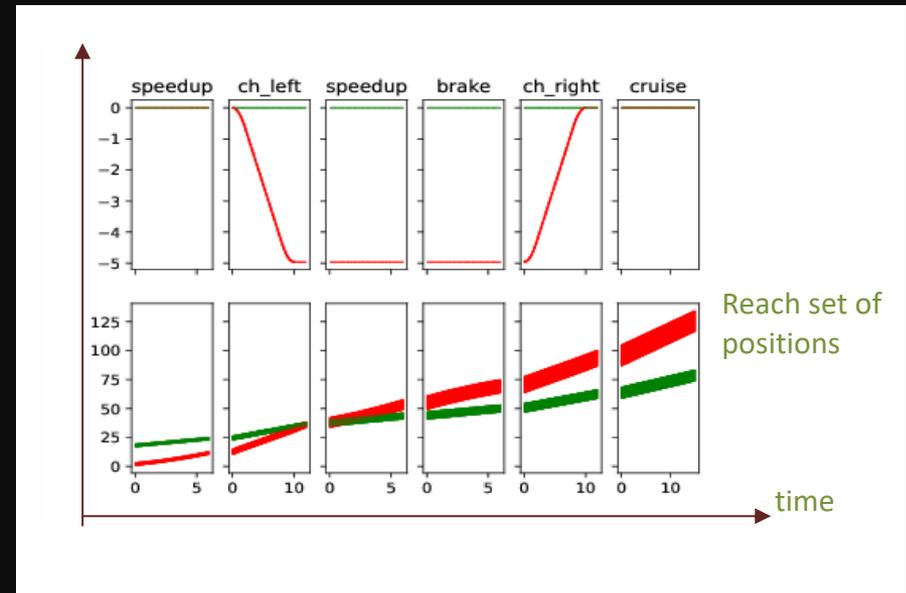
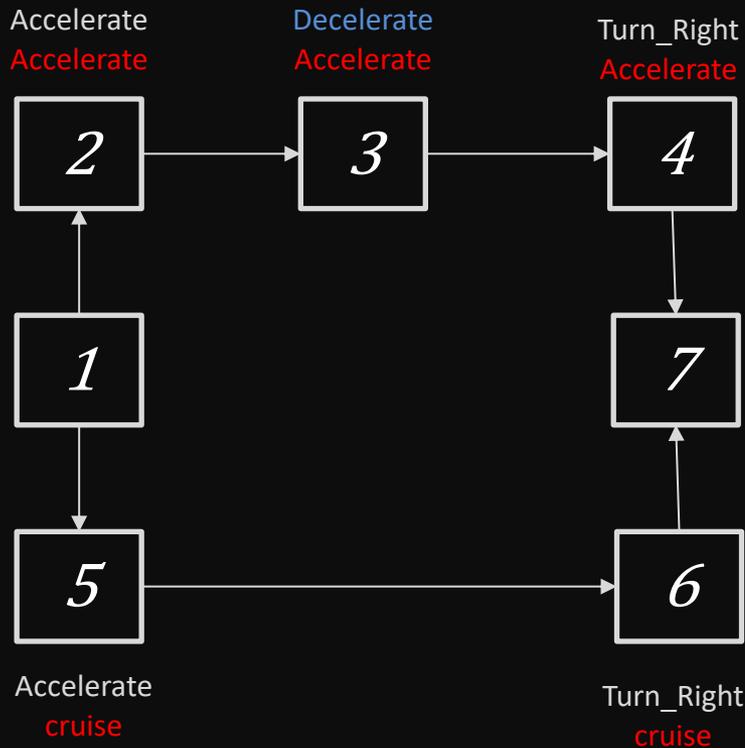


Simulate function takes
as input mode, initial
state, and time horizon

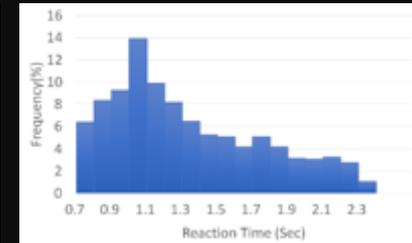
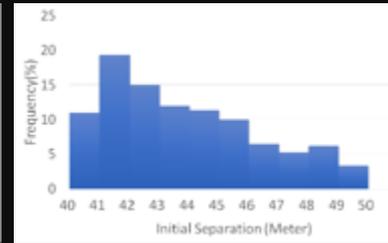
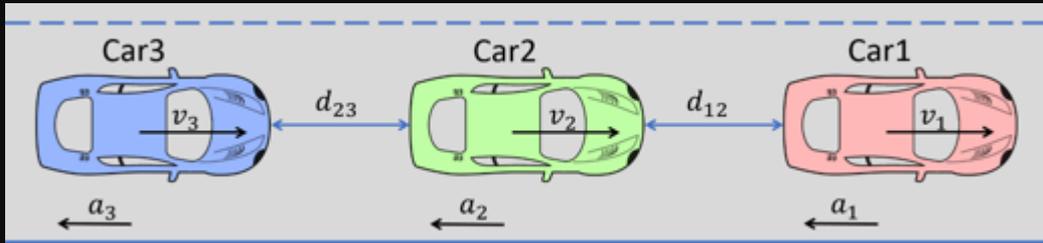
DryVR's Executable
hybrid model



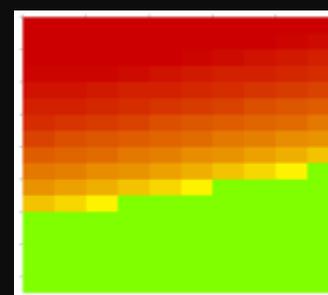
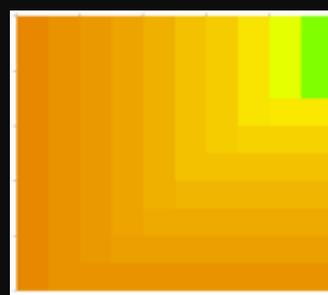
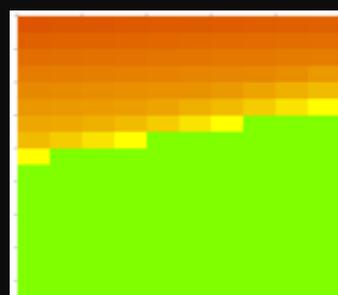
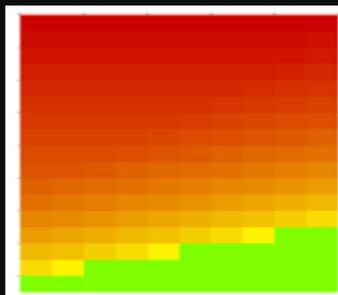
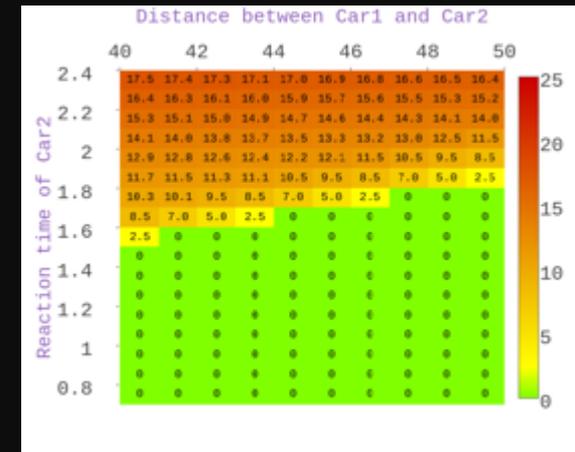
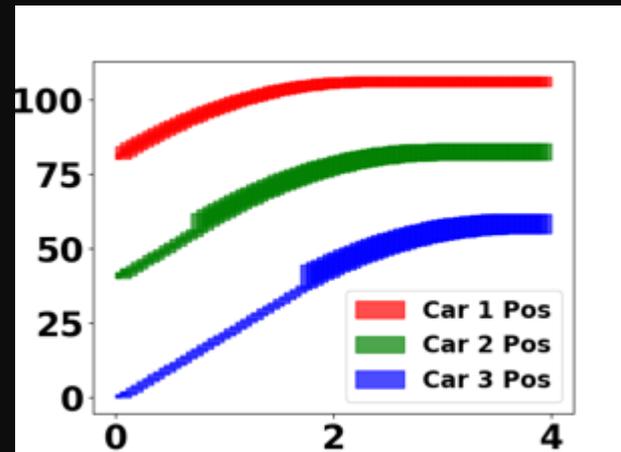
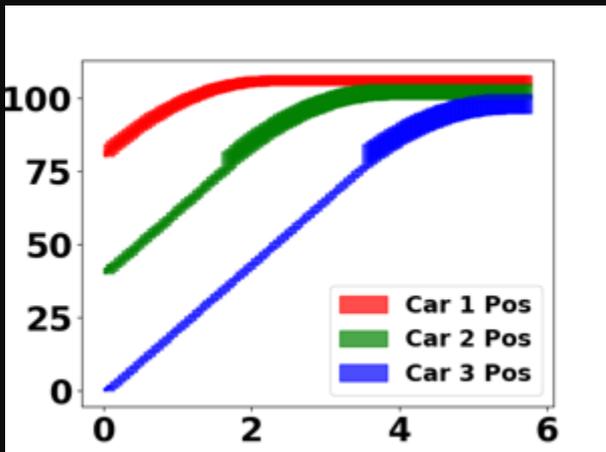
Reachability analysis



Automated Risk / ASIL Analysis



Risk = Probability x Severity



Conclusions

Simulation data + sensitivity from models => algorithms => sound & complete invariance verification

Try C2E2 and DryVR give feedback, built on!

Examples available: Satellites to transistors

Several open questions about handling models with uncertainty and precise characterization of efficiency

This work is supported by grants from the United States National Science Foundation (NSF)